

# Targeted College Admission

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18.05.2025

# Introduction

- Targeted Admission: Matching students, programs, and firms

## Main Contributions:

- 1 Formulate a new market design problem inspired by Russian practice
- 2 Prove impossibility of stable matching between firms, programs, and students
- 3 For fixed firm capacities: Design stable mechanism
- 4 (In progress) Analyze additional properties and alternative frameworks

# Brief History of Targeted Admission in Russia

## Soviet Union:

- 1 1933: Mandatory 5-year job placement for all graduates (Narkom assignments)
- 2 1962: Top workers sent to targeted PhD programs (aspirantura) with return obligations
- 3 1987: Transition to a new system based on TA

## Russian Federation:

- 1 1995: Budget students allowed optional 3-year employment agreements
- 2 2002: Separate quotas for state/municipal employers
- 3 2013-2019: Reserved seats for students with state contracts
- 4 2019: Minimum employment term set to 3 years
- 5 2024: 150000 of total 600000 budget seats were reserved for TA ( 10% - 100% per program)

# TA: Scope and Procedure

## General Admission:

- Students apply to  $\leq 5$  programs across  $\leq 5$  universities
- Rank preferences within each university
- Programs rank applicants by Unified State Exam scores
- Single enrollment consent submitted before deadline
- Universities assign seats via Deferred Acceptance (DA)

## Targeted admission:

- Students select **one** program and **one** employer
- Binding contract for **3–5 years** post-graduation
- TA seats filled before general admission

# TA Shortcomings

**Current inefficiencies stem from regulatory mismatches**

**Student Challenges:**

- Restricted to **one** TA program (strategic uncertainty)
- TA seats allocated first — may lead to suboptimal matches
- Risk of overcommitment to less-preferred options

**Employer Challenges:**

- No explicit capacity controls → potential over-enrollment
- No way to get students committed, thus overbooking
- Risk to get more students

# Literature

## College Admission:

- Gale and Shapley (1962)
- Roth and Sotomayor (1990)

## Affirmative Action in College Admission and School Choice:

- Abdulkadiroğlu and Sönmez (2003)
- Abdulkadiroğlu (2005)

# Literature

## Three-sided matchings:

- Alkan (1987): nonexistence of stable threesome matchings
- Danilov (2003): existence of stable matchings in some three-sided systems (acyclic preferences)
- Boros et al. (2004): stable matchings in three-sided systems with cyclic preferences

# Setting

The targeted college admission problem is defined by the following components:

- A finite set of students  $I = \{i_1, \dots, i_n\}$
- A finite set of schools  $S = \{s_1, \dots, s_m\}$
- A finite set of firms  $F = \{f_1, \dots, f_k\}$
- Each school  $s \in S$  has capacity  $q_s$ , with vector  $q_S = (q_{s_1}, \dots, q_{s_m})$  representing all school capacities
- Each firm  $f \in F$  has capacity  $q_f$ , with vector  $q_F = (q_{f_1}, \dots, q_{f_k})$  representing all firm capacities
- Each student  $i \in I$  has ordinal preferences  $P_i$  over  $S \times F$  (school-firm pairs)



# Setting

- Each school  $s \in S$  has an ordinal priority  $P_s$  over the power set of students  $2^I$ . The priority is assumed to be responsive, meaning that the school's preferences over groups of students are consistent with its preferences over individual students.
- Each firm  $f \in F$  has an ordinal preference  $P_f$  over the set  $2^{S \times 2^I}$ . The preference is also assumed to be responsive.
- The capacity vector  $q = (q_S, q_F)$  combines the capacities of schools and firms.
- The preference profile  $P = (P_I, P_S, P_F)$  represents the preferences of all students, schools, and firms.

The tuple  $(I, S, F, q, P)$  defines the *targeted college admission problem*

# Matching and Mechanism

- A **matching**  $\mu$  is a function
  - Assigns each student to at most one school-firm pair
  - Satisfies  $|\mu(s)| \leq q_s$  for all schools
  - Satisfies  $|\mu(f)| \leq q_f$  for all firms
- $\mu(i)$ : student  $i$ 's assignment ( $\emptyset$  if unmatched)
- $\mu(s)$ : students assigned to school  $s$
- $\mu(f)$ : student-school pairs assigned to firm  $f$
- A **mechanism**  $\phi$  is a function that to each problem finds a matching

Key requirements:

- **Stability**: No blocking by students/firms/schools
- **Respecting Reserves**: No acceptable student is left unmatched while reserves are not filled

# Properties

We identify two key properties

- Preferences are called **homogeneous** if they follow the structure:  $P_{f|s} = P_s$ , meaning that firms rank schools and, within each school, replicate their preferences over students. Preferences may additionally be **sparse** - firms may reject some students, but the ranking order within schools remains unchanged
- A capacity is **fixed** if a firm and a school pre-negotiate the number of slots allocated to students from that school. Otherwise, the capacity is **flexible**: for instance, a firm needs to hire 10 students, but their distribution across programs is unspecified in advance

# Result 1. Impossibility of Stable Matching (Flexible Capacity & Heterogeneous Preferences)

## Setup

$|I| = |S| = |F| = 3, q_{s_i} = q_{f_i} = 1 \forall i \in \{1, 2, 3\}$

Preferences follow *lexicographic relaxation of cyclic preferences*

### Firms' Preferences ( $F$ )

$f_{1,3} : (s_1, i_j) \succ (s_2, i_j) \succ (s_3, i_k)$   
 $f_2 : (s_2, i_j) \succ (s_3, i_j) \succ (s_1, i_k)$   
 $f : (s_j, i_2) \succ (s_j, i_1) \succ (s_j, i_3)$

### Students' Preferences ( $I$ )

$i_{2,3} : (f_1, s_j) \succ (f_2, s_j) \succ (f_3, s_k)$   
 $i_1 : (f_2, s_j) \succ (f_3, s_j) \succ (f_1, s_k)$   
 $i : (f_i, s_2) \succ (f_i, s_1) \succ (f_i, s_3)$

### Schools' Priorities ( $S$ )

$s_{1,3} : (i_1, \cdot) \succ (i_2, \cdot) \succ (i_3, \cdot)$   
 $s_2 : (i_2, \cdot) \succ (i_3, \cdot) \succ (i_1, \cdot)$

## Key Features

- Cyclic preferences create unavoidable blocking pairs
- Flexible capacities prevent stable allocations
- Heterogeneity breaks possible symmetries

## Result 2. Stable Matching under Fixed Capacities

### Key Assumptions

- **Partnership structure:**
  - Either: 1 school  $\leftrightarrow$  1 firm (firms may partner with multiple schools)
  - Or: Treat each school-firm pair as distinct "sub-school" with fixed capacity
- **Fixed capacities:** Pre-negotiated quotas between schools and firms

Under these conditions, a **stable matching exists** and can be found via an adapted Gale-Shapley algorithm.

# Proof Sketch

The algorithm proceeds as follows:

- ① Students apply to top-choice  $(s, f)$  pair
- ② Each school-firm pair  $(s, f)$ :
  - Rank applicants by school's  $P_s$
  - Accept top  $q_{s,f}$  students
  - Reject others
- ③ Rejected students apply to next preference. The process repeats until no further rejections occur

## Result 3. Special Case: Flexible Capacities

### Identical School Rankings

When:

- Schools rank students identically (e.g., via unified exam)

### Dictator Algorithm

- Students processed in exam score order (best to worst)
- Each selects top available school-firm pair
- Yields stable matching by construction

# In progress ...

When capacities are flexible, we hypothesize that stability can be achieved under mild preference restrictions - notably, homogeneous preferences across firms. To prove this, we employ **matching-with-contracts theory**, extending the model to accommodate multi-dimensional constraints while preserving stability conditions.



# Conclusion

Thank you for your attention!

Questions? Discussion?